$\begin{array}{c} {\rm Algorithmics} \\ {\rm Correction~Midterm~\#4~(C4)} \end{array}$

Undergraduate 2^{nd} year (S4) – Epita

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Solution 1 (Cut points, cut edges - 5 points)

- 1. Cut points of G_1 : 1, 7, 8, 10
- 2. Cut edges of G_1 : (1,2), (7,8), (8,9), (10, 11).



Figure 1: Graph, cut points and cut edges of G_1

- 3. The biconnected components of G_1 are :
 - $\{(1,3), (1,6), (1,7), (3,7), (4,6), (4,7), (5,6), (5,7)\}$
 - $\{7,10\}, (7,12), (10,12), (10,13), (12,13)\}$
 - {(1,2)}
 - {(7,8)}
 - {(8,9)}
 - {(10,11)}



Figure 2: 2-connected components of G_1

4. The table of *prefix* and *higher* values is :

	prefix	higher
1	1	1
2	2	2
3	3	1
4	5	1
5	7	4
6	6	1
7	4	1
8	8	8
9	9	9
10	10	4
11	11	11
12	12	4
13	13	10



Figure 3: Table of values obtained during the depth-first traversal of ${\cal G}_1$

Figure 4: spanning forest associated to the depth-first traversal of ${\cal G}_1$

Solution 2 (I want to be a tree -8 points)

- 1. Definitions:
 - \Box A *tree* is an acyclic connected graph.
 - \Box A tree is a connected graph with n-1 edges (n: vertex number).
 - \Box A tree is an acyclic graph with n-1 edges (n: vertex number).
 - \Box A *tree* is an acyclic graph and adding any edge creates a cycle.
 - \Box A *tree* is a connected graph that is no longer connected after the removal of any vertex.
- 2. (a) Edges that can be removed: Those who become back edges.
 - (b) the list of the edges of the graph "Not a tree yet" removed:

11 - 9 12 - 8 5 - 2 10 - 4

- 3. During the depth-first search, we assign to each vertex the number of the connected component it belongs to (from 1 to k, if there are k components):
 - (a) Number of edges to add: k-1
 - (b) What are the edges to add, during the traversal? For example, you can add edges from the first vertex chosen for the traversal to all the roots of the other trees.
 - (c) the list of the edges of the graph "Not a tree yet" added:

For instance 0 - 1 and 1 - 2

4. Specifications:

The function $make_me_tree(G)$ turns the graph G into a tree and returns the connected component vector of the initial graph.

```
def __makeMeTree(G, s, p, cc, noc):
1
           cc[s] = noc
2
           for adj in G.adjlist[s]:
3
               if cc[adj] == 0:
4
                    __makeMeTree(G, adj, s, cc, noc)
5
               else:
6
                    if adj != p:
7
                        G.removeedge(s, adj)
8
9
      def makeMeTree(G):
10
           cc = [0] * G.order
           x = 0
12
           noc = 1
13
           __makeMeTree(G, 0, -1, cc, noc)
14
15
           for y in range(1, G.order):
               if cc[y] == 0:
16
                    noc += 1
17
                    __makeMeTree(G, y, -1, cc, noc)
18
                    G.addedge(x, y)
19
20
                    x = y
           return cc
```

Solution 3 (Condensation – 4 points)

Specifications:

The function condensation(G, scc) builds the condensation G_R of a digraph G, with scc its strongly connected component list. The function returns G_r and the vector of strongly connected components: a vector that gives for each vertex the number of the component it belongs to (the vertex in G_R).

```
def condense(G, scc):
2
3
         comp = [-1] * G.order
4
         k = len(scc)
         for i in range(k):
5
             L = scc[i]
                                                 for s in scc[i]:
6
                                                     comp[s] = i
             for j in range(len(L)):
7
                  comp[L[j]] = i
8
9
         Gr = graph.Graph(k, directed = True)
         for s in range(G.order):
             for adj in G.adjlists[s]:
12
                  (x, y) = (comp[s], comp[adj])
13
14
                  if x != y and y not in Gr.adjlists[x]:
                      Gr.addedge(x, y)
                                            \# Gr. a d j l i s t s [x]. append (y)
16
17
         return (Gr, comp)
```

Solution 4 (Graphes and Mystery – 3 points)

1.

	Call number	Returned result
(a) test(G_2)	5	False
(b) test(G_3)	7	True

2. What is the information returned by test(G)?

test(G) checks if G is strongly connected.

Solution 5 (Saving Algernon – Bonus)

- 1. (a) The kidnapper is the one of the lab # 1.
 - (b) Algernon is in the air vent j.
- 2. We search for an Eulerian path in the graph below, where each zone is a vertex (between 2 detectors). The only two vertices of odd degrees are A and T. A is the only zone that contains an access to a laboratory (the # 1): this is Algernon's starting point. Thus H is the end point: Algernon is hidden in the air vent j.

